

## Editorial: Horses

To solve 1st subtask we need just calculate our answer using dynamic programming  $dp[i][j]$  - what is maximal profit if we pass  $i-1$  days and we have  $j$  horses then we got  $j * x_i$  horses and check all number of horses that we sell today and go to the next day.

### #1 Observation

First of all let's consider to points  $i$  and  $j$ , what if we sell  $k$  horses in  $i$ -th and  $j$ -th day. And check in what day it's more preferable.

Profit of  $i$ -th day:  $x_1 * x_2 * \dots * x_i * y_i$

Profit of  $j$ -th day:  $x_1 * x_2 * \dots * x_i * x_{i+1} * \dots * x_j * y_j$

It depends on this

$y_i ? x_{i+1} * \dots * x_j * y_j$

>

<

=

so if it's > then it's profitably sell horses in  $i$ -th day

if it's < then it's profitably sell horses in  $j$ -th day

if it's = then it's no matter when we sell them in  $i$ -th day or in  $j$ -th day

From this point it's clear that it is better to sell all our horses in one day that gives us maximal profit.

Using this observation we can solve 2 subtasks.

After each query we can solve problem in  $O(n)$

### #2 Observation

If all  $x_i \geq 2$ , then after 30 days  $x_i * x_{i+1} * x_{i+2} * \dots * x_{i+29} \geq 2^{30} > 10^9$ , so position less than and equal  $n-30$  never can be optimal solution, because even if  $y_{n-30} = 10^9$  and  $y_n = 1$ ,  $2^{30} * y_n > y_{n-30}$

It's enough to check only last 30 positions

### #3 Observation

The main problem in last subtask is  $x_i = 1$ , but in that cases multiplication first  $x_i$  numbers and  $x_{i-1}$  is not change, this gives us opportunity to merge consecutive positions with  $x_i = 1$ , when we merge this positions we should take the maximal  $y_i$ . So if we do that, it's enough to check last 60 positions, because it can be no more than 30 merged 1's between 30 last positions where  $x_i \geq 2$ .

So we can store our state in some structure like set, that provide us information about current merged positions, and some structure that provide us RMQ. So solution per query will be  $\log(10^9) * \log(n)$ . Then we have to calculate answer we can use something like segment tree.

So overall complexity will be  $O(m * \log(10^9) * \log(n))$ .